

Assessment of Congestion Control Technique in Wireless Mesh Network Using Routing Protocols

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Abstract

In Multi-Radio, Multi-Channel, remote Mesh Networks we address the issue of congestion control. This issue is altogether more intricate on the grounds that it obliges the channel alterations and the movement allotments for every channel be mutually advanced. The path in which the blockage is taken care of may bring about more postpone and more bundle misfortune and a huge overhead may happen. We address these issues by directing method by utilizing definition that permits its disintegration in blockage control over a hub way. We look at some directing methods that achieve high throughput and build in parcel conveyance degree and lessened end to end delay.

Keywords

Protocol, Control, WME, Congestion, Systems.

I. Introduction

Wireless mesh networks (WMNS) are alterably dealt with toward oneself and self-arranged, with the hubs in the system consequently creating a specially appointed system and keeping up the cross section connectivity's lattice system is a lattice topology in which every hub transfers information for the system. A cross section arranges whose hubs are all joined with one another is completely associated systems. There exists more than one way between a source and an objective in a system. It obliges conventions in WMNS to be distributive and communitarian. WMNS still have exceptionally restricted capacities of incorporating heterogeneous remote systems, because of the trouble in building different remote interfaces and the relating portal/scaffold works in the same cross section switch In a system when the amount of bundles sent to any system is more than it can deal with, blockage is produced. Elements like parcel landing rate surpassing the friendly connection limit, inadequate memory to store arriving bundles cause congestion. It's a worldwide issue including each switch and have inside the sub system blockage control is concerned with the productivity of the system at high load. To diminish the congestion in a WMN's some directing conventions are utilized to make it more proficient and solid. Steering conventions for specially appointed systems could be ordered into two fundamental classifications: Proactive Driven directing conventions and On-interest steering conventions. Correlation between the steering conventions focused around Packet conveyance rate (PDR), End to End Delay and Congestion. This paper means to give a similar investigation of 3 well known steering conventions AODV, DSR, AOMDV and DSDV for Multi radio Wireless Mesh Networks. Whatever is left of the paper is sorted out as takes after: Section 2 introduces an outline of the Wireless steering convention that is examined and looked at. Area 3 gives a short depiction of the Simulation parameters, suspicions hold and portrayal of the orderly looking at approach utilized within the paper. Area 4 provides the reproduction comes about and examines it. At last the conclusion is given in area 5.

II. Background Work

Portable Ad Hoc Networks could be isolated into Table-Driven and On-Demand Routing convention where table Driven conventions

are proactive and keep up a steering table and On-Demand are dynamic and don't keep up a directing table. The accompanying directing conventions are examined in the examination:

A. Destination-Sequenced Distance Vector routing (DSDV)

DSDV is a table driven directing plan for impromptu portable systems focused around the Bellman-passage calculation. The change made to the Bellman-Ford calculation incorporates flexibility from circles in directing table by utilizing arrangement numbers. Every hub demonstrations as a switch where a directing table is kept up and intermittent steering overhauls are trade, regardless of the fact that the courses are not required. A succession number is connected with each one course or way to the terminus to avoid steering circles. Steering redesigns are traded regardless of the possibility that the system is unmoving which uses up battery and system transmission capacity. Therefore, it is not best for very dynamic systems.

B. Ad hoc On-Demand Distance Vector Routing (AODV)

AODV is an on-demand steering convention which is juncture of DSDV and DSR. Course is ascertained on interest, in the same way that it is in DSR through course disclosure process. Be that as it may, AODV keeps up a steering table where it keeps up one passage for every goal dissimilar to the DSR that keeps up numerous course reserve entrances for every terminus. AODV gives circle free courses while repairing connection breakages however not at all like DSDV, it doesn't require worldwide intermittent directing notices.

C. Ad hoc On-Demand Multipath Distance Vector Routing (AOMDV)

Impromptu On-interest Multipath Distance Vector Routing (AOMDV) convention is a growth to the AODV convention for figuring different circle free and connection disjoint ways. The directing sections for every end of the line hold an arrangement of the following bounces alongside the relating Hop checks. All the following jumps have the same arrangement number. This aides in staying informed concerning a course. For every end of the line, a hub keeps up the promoted Hop tally, which is

characterized as the greatest bounce mean all the ways, which is utilized for sending course ads of the destination in DSR, the entire course is conveyed with the message as an overhead, inasmuch as in AODV, the directing table is kept up consequently it is not needed to send the entire course with the message throughout the Route Discovery process

D. Dynamic Source Routing (DSR)

Element Source Routing is a Pure On-Demand directing convention, where the course is ascertained just when it is needed. It is intended for utilization in multi bounce impromptu systems of portable nodes. dsr permits the system to act naturally sorted out and self-arranged without any focal organization and system framework. It utilizes no occasional steering messages like AODV, subsequently decreases data transmission overhead and preserved battery control and additionally vast directing overhauls. It just needs the exertion from the MAC layer to recognize join failure. dsr utilization source steering where the entire course is conveyed as an overhead.

III. Methodology Implemented

We are making a system of 50 hubs comprising of static and half and half in a framework topology. This system comprises of multi-source and multi objectives to contrast the directing conventions and itself and different conventions. The dissection is performed utilizing Network Simulator (Ns-2), which is prevalently utilized for impromptu systems administration group. The steering conventions were analyzed focused around the 4 performance measurements:

Packet Delivery Rate (PDR): The ratio of Data packets delivered to those generated by the sources.

End to End delay: the delay in delivering a packet to the destination which is inclusive of all kinds of delay.

Normalized Routing Overload: This is the routing packets sent per Delivered packet at the destination.

Throughput: The rate of successful message delivery over a communication channel.

The transmission of information is performed utilizing a convention called TCP, which is an association arranged protocol. tcp, is utilized for high solid systems so the transmission time is generally less. tcp convention likewise has an effective utilization for congestion control.

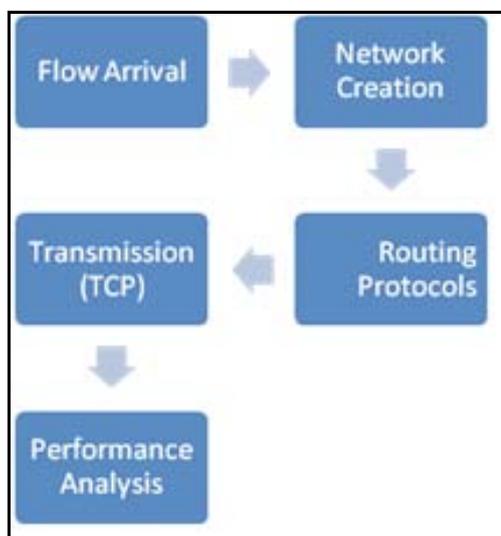


Fig. 1: Process Flow

A. Network Architecture

In this building design, network switches structure a base for customers, as indicated in Fig. 1, where dashed and strong lines show remote and wired links, respectively. The WMN base/spine might be manufactured. Utilizing different sorts of radio advances, notwithstanding the basically utilized IEEE 802.11 innovations. The cross section switches structure a lattice of masterminding toward oneself patching toward oneself connections among themselves. With passage usefulness, network switches might be connected to the Internet. This methodology, additionally alluded to as infrastructure fitting,

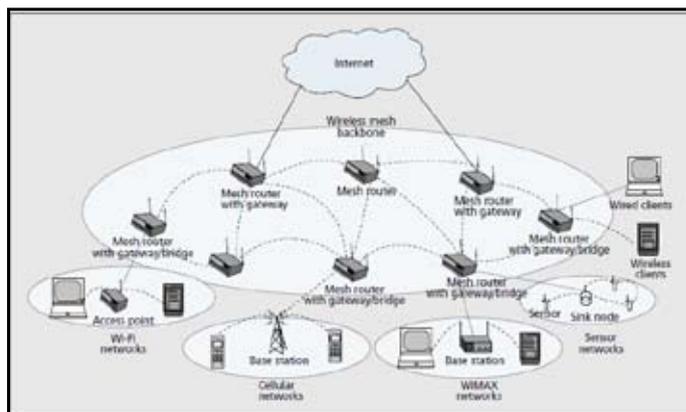


Fig. 2: Infrastructure WMN's

The system secured consists of static and half and half hubs which comprise of multisource-terminus sets. Three situations of one source one terminus, three sources three ends and six sources and six goals are secured. For every situation the information transmission from the source and the end is demonstrated in the beneath figures.

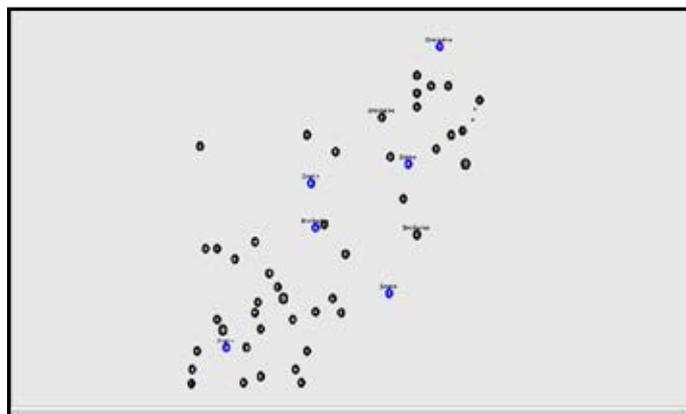


Fig.3: Sample network with 6 source and 6 destinations

B. Simulation Method

Before we begin the reproduction, we make 3 format TCL scripts to be utilized by our group record to automatically reproduce situations utilizing the Mobility scene produced by utilizing the set dest toolset. 4 Batch documents are used: batch record to run the reproductions focused around the test situations changing speed and stop time, group document to duplicate the test situations in the format tell script, the bunch document to run the awk script and the last cluster document to move the name, follow and versatility situations in particular envelope for filing and future utilization. The reproductions parameters are indicated in the table. The same set of versatility situations for every variety of pace and stop time

is utilized while changing the directing protocol.

Studied protocols	: DSDV, AODV, AOMDV and DSR.
Simulation time	: 1000 seconds
Simulation area	: 500 m X 500 m
Node Movement model	: Random Waypoint
Speed	: 0-25 in steps of 5 m/s
Traffic type	: CBR(UDP)
Node Pause time	: 0-100s in steps of 20s
Data payload	: 1000 bytes/ Packet
Packet Rate	: 250 k
No. of Nodes	: 50
No. of source destination pairs	: 1,3,6

Fig. 4: Tabulation of Parameters

IV. Results

Recreations were carried out fluctuating the velocity keeping the stop time steady (0 sec) and afterward changing the Stop time keeping the rate consistent (5 m/s).the variety were carried out separately for all the steering Convention from AODV to DSR to DSDV and AOMDV. The amount of sources and objectives for correlation were likewise shifted from 1 to 3 to 6 to distinguish the impact. In all situations the Comparison were focused around execution metric: Packetdelivery Ratio, end to end Delay, Throughput and standardized Routing Load.

A. Comparison based on End to End Delay

As it might be seen from the above results, end to end postponement is higher in DSR took after by AODV, AOMDV and DSDV having the least and most steady End to End Delay. DSR is an On-Demand source directing convention, and this is the significant explanation behind it having a higher End-to-End Delay, where course is looked just when required and there is a course Discovery component happening each time and it likewise need to convey a huge overhead each one time, accordingly the higher deferral. AODV then again has one and only course for every end in the directing table, which is always redesigned focused around grouping number and DSDV need to constantly redesign the entire steering table intermittently and when required, which prompts a slight postpone in conveyance. AOMDV has numerous courses for every objective which stays stable therefore the postponement is likewise steady. The end to end delay does change with increment in the amount of source –destination matches as the source and objective moving with variety of pace, the expanded number of source-goal combines just may expand number of bounces. The End to End deferral abatements with expansion with rate, as when it moves all the more much of the time the steering upgrades are traded all the more habitually and speedier it achieves the terminus. It is apparent that the End to End Delay increments with expansion in number of source-objective combines, so does the amount of bounces and along these lines the delay.dsr has a stable postpone and is more dependable took after by DSDV and AODV and AOMD

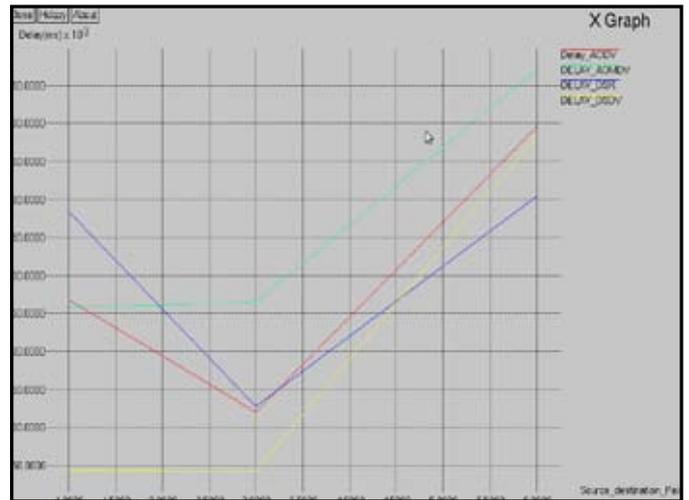


Fig. 5: End to End Delaycomparison of the Routing protocols

B. Comparison based on Congestion

As it could be seen and expected the Routing burden is higher in DSR than in DSDV and AODV, as DSR is asource steering convention and the entire course is sent with each one message as an overhead, which builds the directing load.it increments with pace also, as all the more much of the time the end of the line changes the more course is conveyed as an overhead. It is steady on account of AODV and DSDV as they keep up a directing table and have fewer courses overhead



Fig. 6: Congestion comparison of the Routing protocols

C. Comparison based on PDR

As it can be seen from the Fig 7, the PDR is the same for AODV,AOMDV and DSR followed by DSDV which has the least PDR. In the scenario of increasing the source-destination pair which has the multi hop characteristics, AOMDV has slightly more PDR than AODV and DSDV .DSR has slightly more PDR than AODV as it always looks for the most fresh and reliable route when needed and does not look for it from the routing table like AODV.

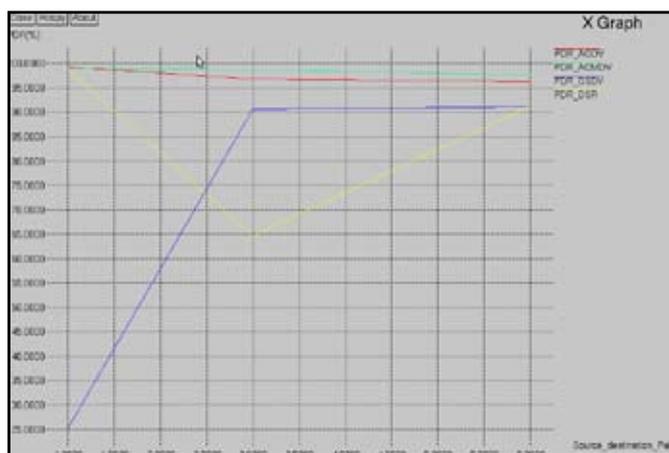


Fig. 7: PDR comparison of the Routing protocols

Conclusion

This paper looked at the 4 famous steering conventions AODV, Dsr,aomdv and DSDV .Simulation results demonstrate that amongst all the conventions, AODV has a stable End to End Delay notwithstanding versatility as it has the gimmick of On-Demand Routing convention. AOMDV has a higher PDR than the other three steering conventions in portability as it is a Table Driven convention and is more reliable. dsr has the least End to End Delay and blockage. In view of the above reproduction situations, parameter, supposition and results AODV could be considered as an effective speedier directing convention than DSR and DSDV however has a higher Routing load similarly than DSDV. DSR is albeit more steady and profoundly dependable.

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